

EXTENDING THE COMPUTATIONAL HORIZON

Attila Egri-Nagy

Centre for Research and Mathematics
School of Computing, Engineering and Mathematics
University of Western Sydney, Australia

RESEARCH QUESTIONS

- **Bacterial Genomics** (with Andrew Francis and Volker Gebhardt)
 - How to reconstruct phylogeny trees?
- **Foundations of Computing** (with James East, James D. Mitchell, Chrystopher L. Nehaniv).
 - What is computable with n states?
 - What is the structure of finite computations?

Precise answers can be obtained in *abstract algebra*, in computational group and semigroup theory.

WHY COMPUTATIONAL?

- More examples, more raw data for the mathematical reasoning.

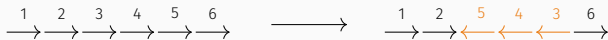
BACTERIAL CHANGES

Single celled organisms with circular chromosome.

- *Local* changes such as *single nucleotide polymorphisms* (SNPs):

ACGGCCCTTAGG → ACGGCCATTAGG

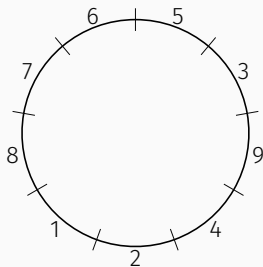
- *Regional* changes such as *inversion* that affect whole regions along the chromosome.



(Regional changes include inversion, translocation, deletion and others.)

- *Topological* changes that produce knots and links in the DNA.

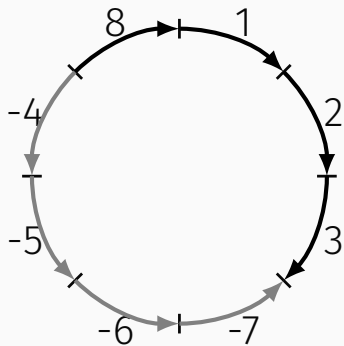
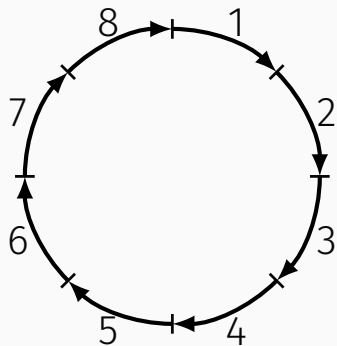
Genome → permutations



Sequences of evolutionary → Sequences of generators

Genomic distance → Length of geodesic words

Genomic space → Cayley-graph



Reference genome and the signed permutation
 $[1, 2, 3, -7, -6, -5, -4, 8]$.

WHERE DOES THE DIFFICULTY FROM?

- The groups are finite and well-studied (symmetric, hyperoctahedral), but big, e.g. \mathcal{S}_{80}
- The generating sets are unusual, “biological”.

For example,

- 2-inversions of the circular genome (vs. linear)
- looking at the “width” as well

Strategy: Calculate and look at small, but non-trivial examples to get insights.

SYMMETRIES OF IRREDUCIBLE GENERATING SETS

	1	\mathbb{Z}_2	$\mathbb{Z}_2 \times \mathbb{Z}_2$	\mathbb{Z}_3	D_8	\mathcal{S}_3	\mathcal{S}_4	\mathcal{S}_5
\mathcal{S}_2		1						
\mathcal{S}_3	1	1						
\mathcal{S}_4	8	5				1		
\mathcal{S}_5	150	25		1		1	1	
\mathcal{S}_6	7931	645	11	6	4	20	2	2

Definition

A *semigroup* is a set S with an associative binary operation $S \times S \rightarrow S$.

Example (Flip-flop monoid)

	1	a	b
1	1	a	b
a	a	a	b
b	b	a	b

ABSTRACT AND TRANSFORMATION SEMIGROUPS

Definition

A *semigroup* is a set S with an associative binary operation $S \times S \rightarrow S$.

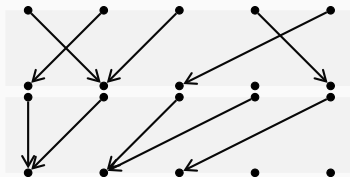
Example (Flip-flop monoid)

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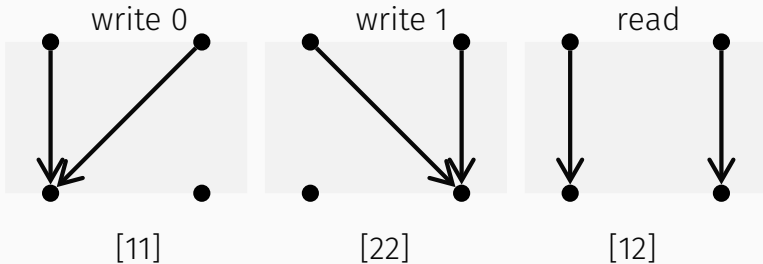
Definition

A *transformation semigroup* (X, S) is a set of states X and a set S of transformations $s : X \rightarrow X$ closed under function composition.

Example (Transformations)



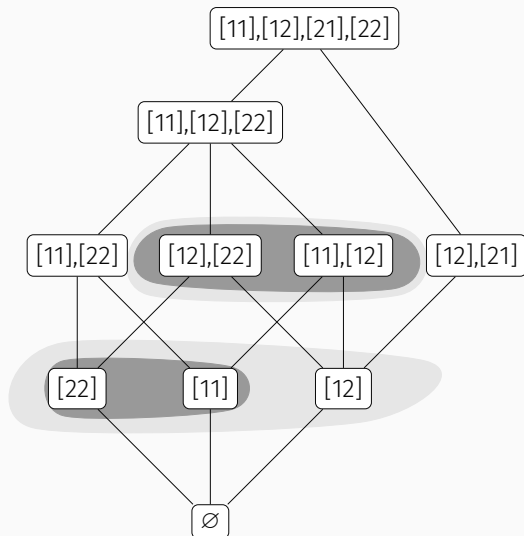
FLIP-FLOP, THE 1-BIT MEMORY SEMIGROUP



So these are computational devices... \approx automata

With transformation semigroups, we get all semigroups.
(Cayley's theorem)

DEGREE 2 TRANSFORMATION SEMIGROUPS

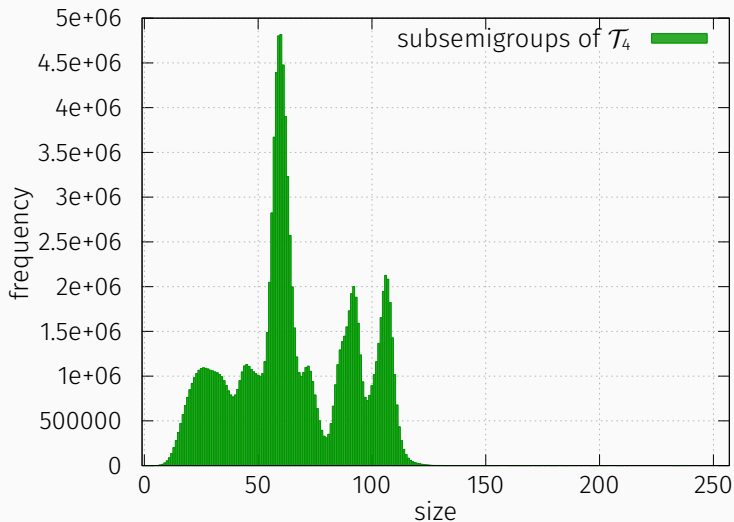


Number of subsemigroups of full transformation semigroups.

	#subsemigroups	#conjugacy classes	#isomorphism classes
\mathcal{T}_0	1	1	1
\mathcal{T}_1	2	2	2
\mathcal{T}_2	10	8	7
\mathcal{T}_3	1 299	283	267
\mathcal{T}_4	3 161 965 550	132 069 776	131 852 491

After discounting the state-relabelling symmetries the database of degree 4 transformation semigroups is still around 9GB.

SIZE DISTRIBUTION



SIZE DISTRIBUTION – LOGARITHMIC SCALE

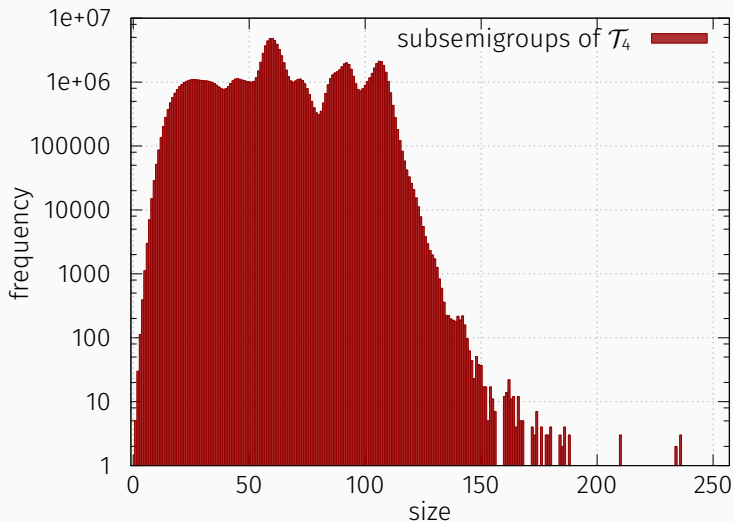
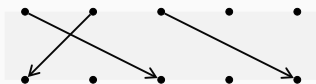


DIAGRAM SEMIGROUPS – TYPICAL ELEMENTS



DIAGRAM SEMIGROUPS – TYPICAL ELEMENTS



$\in \mathcal{I}_n,$



$\in \mathcal{B}_n$



$\in \mathcal{T}_n,$



$\in \text{TL}_n$



$\in \mathcal{S}_n,$



1_n

DIAGRAM SEMIGROUPS

$$\mathcal{P}_1 \hookrightarrow \mathcal{T}_2$$

$$\mathcal{P}_2 \hookrightarrow \mathcal{T}_5$$

$$\mathfrak{B}_1 \cong \mathcal{T}_1$$

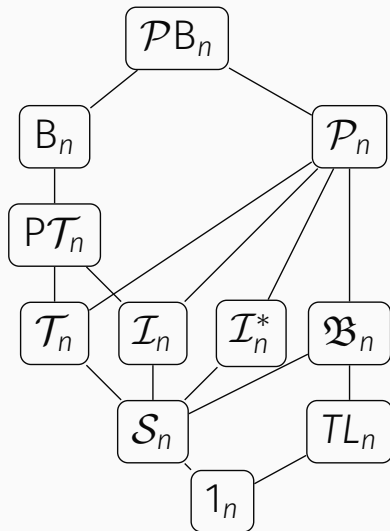
$$\mathfrak{B}_2 \hookrightarrow \mathcal{T}_3$$

$$\text{TL}_1 \cong \mathcal{T}_1$$

$$\text{TL}_2 \hookrightarrow \mathcal{T}_2$$

$$\text{TL}_3 \hookrightarrow \mathcal{T}_4$$

$$\mathcal{P}_1 \hookrightarrow \mathfrak{B}_2$$



	Order	$n = 1$	$n = 2$	$n = 3$	$n = 4$	$n = 5$	$n = 6$
\mathcal{PB}_n	$2^{(2n)^2}$	16	65536	2^{36}	2^{64}	2^{100}	2^{144}
B_n	2^{n^2}	2	16	512	65536	2^{25}	2^{36}
\mathcal{P}_n	$B_{2n} = \sum_1^{2n} S(2n, k)$	2	15	203	4140	115975	4213597
\mathcal{PT}_n	$(n+1)^n$	2	9	64	625	7776	117649
\mathcal{I}_n^*	$\sum_1^n k! (S(n, k))^2$	1	3	25	339	6721	179643
\mathcal{T}_n	n^n	1	4	27	256	3125	46656
\mathcal{I}_n	$\sum_0^n k! \binom{n}{k}^2$	2	7	34	209	1546	13327
\mathfrak{B}_n	$(2n-1)!!$	1	3	15	105	945	10395
\mathcal{S}_n	$n!$	1	2	6	24	120	720
TL_n, J_n	$C_n = \frac{1}{n+1} \binom{2n}{n}$	1	2	5	14	42	132

COMPUTATIONAL HORIZON

	$n = 1$	$n = 2$	$n = 3$	$n = 4$	$n = 5$	$n = 6$
\mathcal{PB}_n	1262					
B_n	4	385				
\mathcal{P}_n	4	272				
\mathcal{PT}_n	4	50	94232			
\mathcal{I}_n	4	23	2963			
\mathcal{I}_n^*	2	6	795			
\mathcal{T}_n	2	8	283	132069776		
\mathcal{B}_n	2	6	42	10411		
\mathcal{TL}_n	2	4	12	232	12592	324835618
\mathcal{S}_n	1	2	4	11	19	56

Given a semigroup S , the equivalence relation \mathcal{J} is defined by

$$t \mathcal{J} s \iff S^1 t S^1 = S^1 s S^1,$$

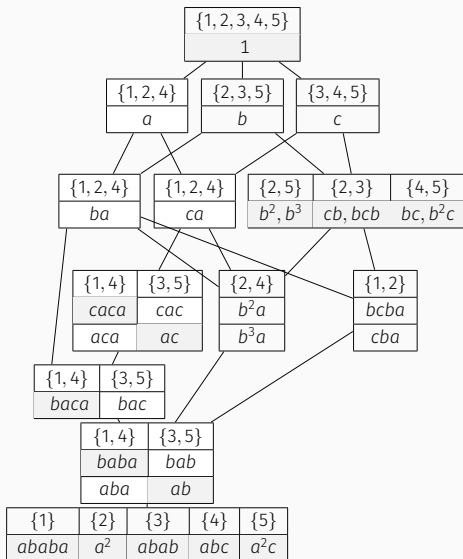
where S^1 is S with an identity adjoined in case S is not a monoid.

In other words,

$$t \mathcal{J} s \iff \exists p, q, u, v \in S^1 \text{ such that } t = psq \text{ and } s = utv$$

The equivalence classes of \mathcal{J} are “*local pools of reversibility*”.

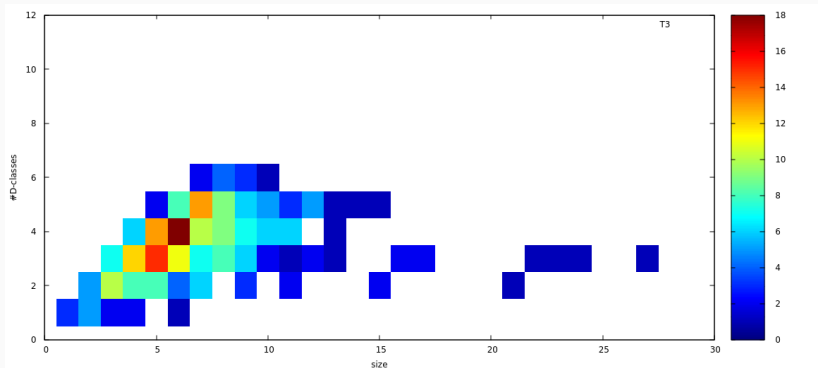
$a = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 2 & 2 & 1 & 2 & 4 \end{pmatrix}$, $b = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 3 & 5 & 2 & 3 & 2 \end{pmatrix}$, $c = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 3 & 5 & 4 & 5 & 4 \end{pmatrix}$ and $M = \langle a, b, c \rangle$. $|M| = 31$



TRANSFORMATION SEMIGROUPS OF DEGREE 3

x axis : size of the semigroups

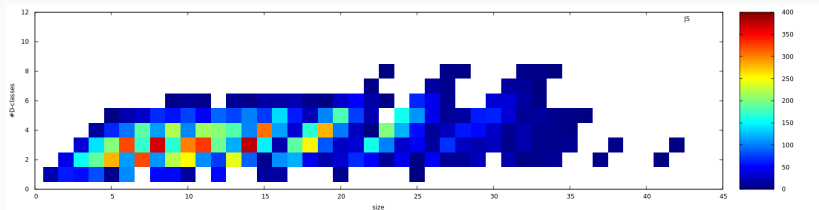
y axis : the number of \mathcal{D} -classes



SUBSEMIGROUPS OF THE DEGREE 5 JONES MONOID

x axis : size of the semigroups

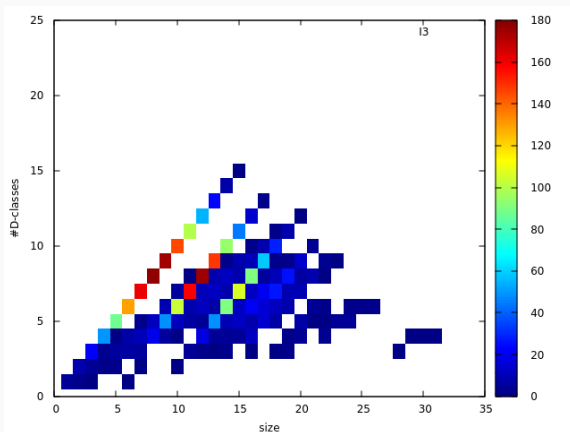
y axis : the number of \mathcal{D} -classes



INVERSE SEMIGROUPS (OF PARTIAL PERMUTATIONS)

x axis : size of the semigroups

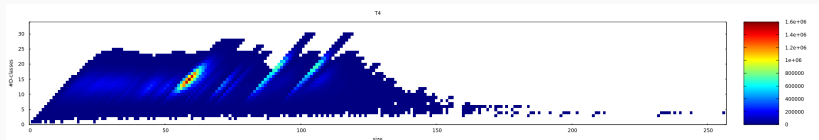
y axis : the number of \mathcal{D} -classes



TRANSFORMATION SEMIGROUPS OF DEGREE 4

x axis : size of the semigroups

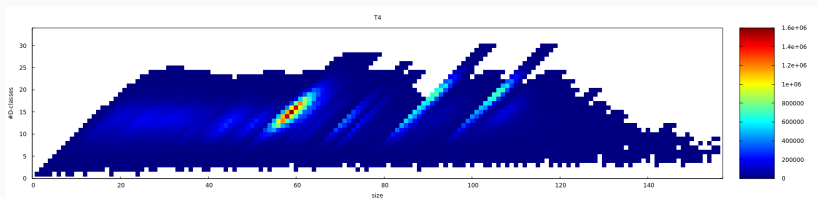
y axis : the number of \mathcal{D} -classes



TRANSFORMATION SEMIGROUPS OF DEGREE 4

x axis : size of the semigroups

y axis : the number of \mathcal{D} -classes



EVEN MORE STRUCTURE, ALL GREEN'S RELATIONS

\mathcal{L} , \mathcal{R} equivalence relations

$$t \mathcal{R} s \iff tS^1 = sS^1,$$

$$t \mathcal{L} s \iff S^1t = S^1s,$$

$$t \mathcal{J} s \iff S^1tS^1 = S^1sS^1$$

$$\mathcal{L} \circ \mathcal{R} = \mathcal{R} \circ \mathcal{L} = \mathcal{D}$$

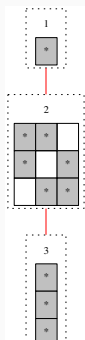
$\mathcal{J} = \mathcal{D}$ in the finite case

$$t \mathcal{R} s \iff \exists p, q \in S^1 \text{ such that } t = sp \text{ and } s = tq$$

$$t \mathcal{L} s \iff \exists p, q \in S^1 \text{ such that } t = ps \text{ and } s = qt$$

“EGGBOX” PICTURE

Tables are \mathcal{D} -classes. Columns are \mathcal{L} -classes, rows are \mathcal{R} -classes. Shaded cells are \mathcal{H} -classes that contain idempotents – used for locating subgroups of the semigroup. \mathcal{T}_3 :



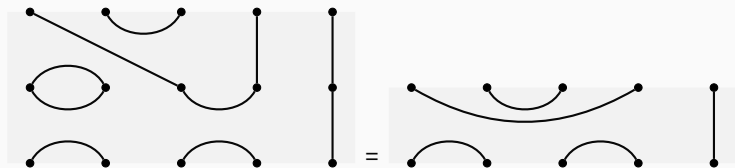
TEMPERLEY-LIEB, JONES MONOID

Catalan numbers, sequences of well-formed parentheses.



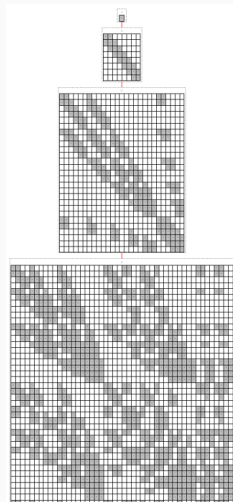
corresponds to $((()()))()$

Applications in Physics: statistical mechanics, percolation problem.



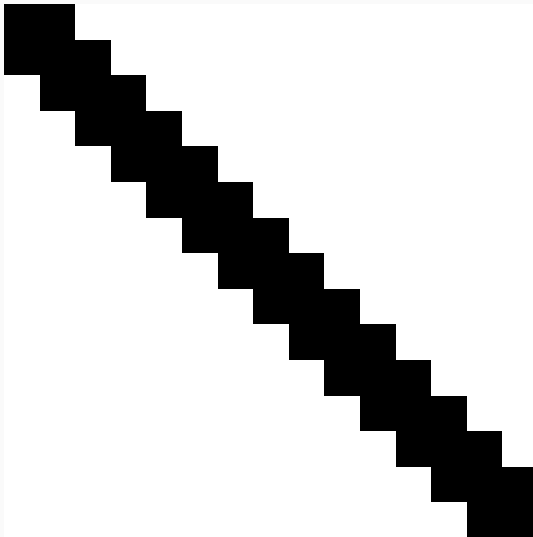
FIRST GLIMPSE THROUGH THE USUAL EGGBOX DIAGRAMS

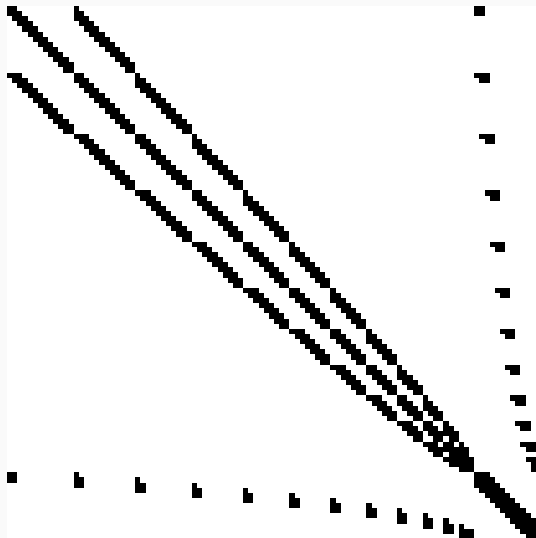
J_9



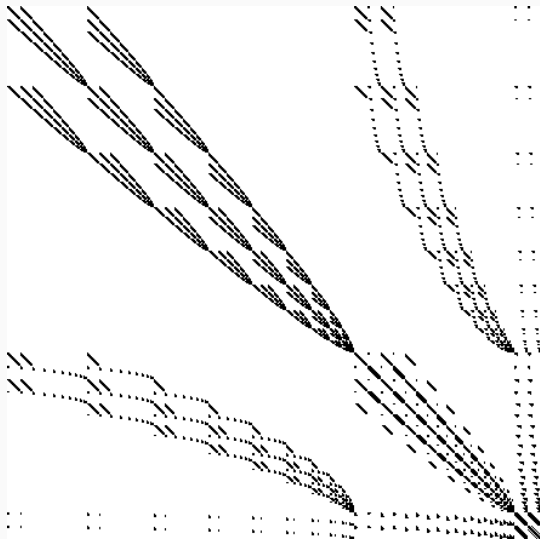
J_{16}

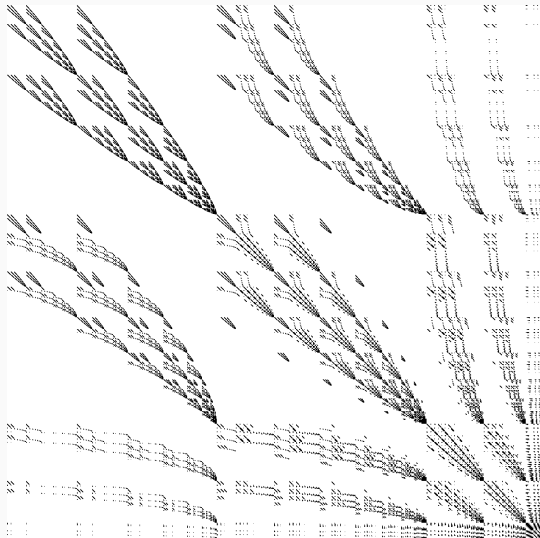


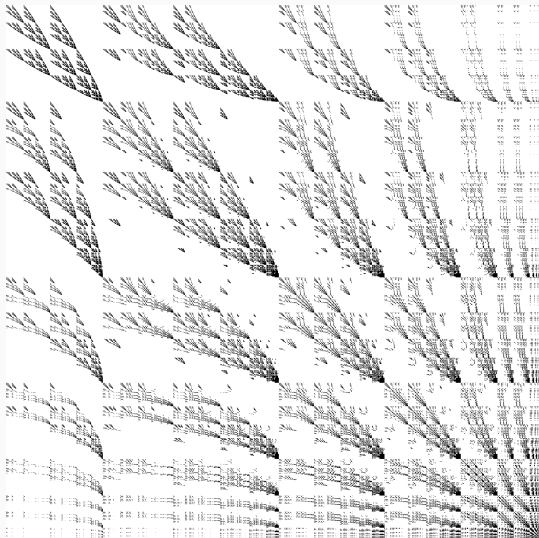


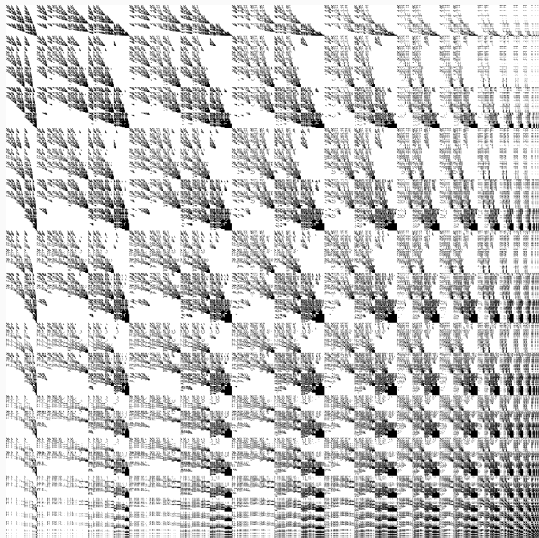


LEVEL 4

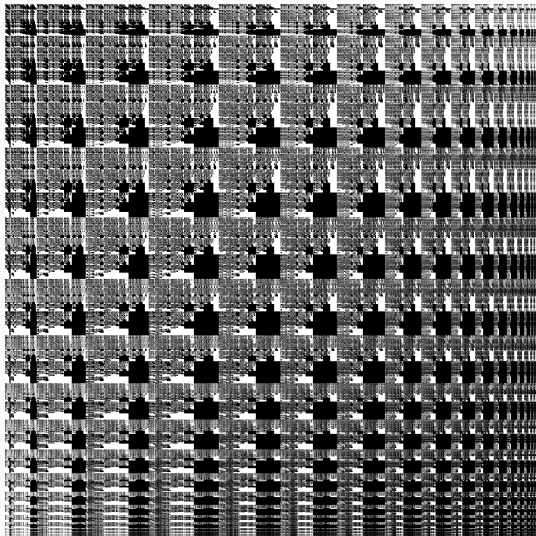


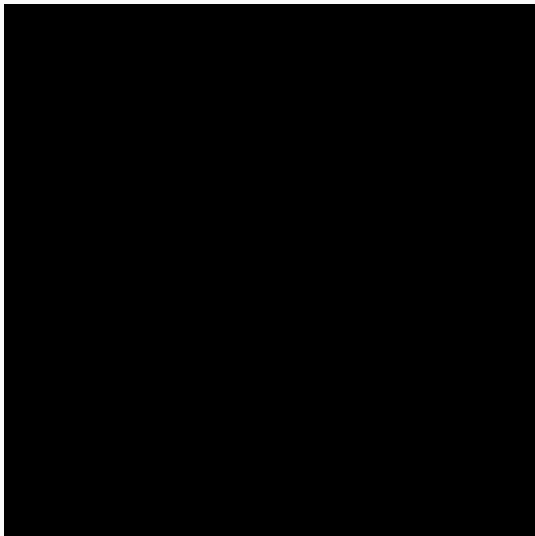






LEVEL 8





The Good We can discover/construct more and more new, interesting and useful mathematics by using computers.

The Bad There is a gap between mathematical rigour and the correctness of software implementations and the physicality of computation.

and The Ugly Developing software is still detrimental to academic career.

Blog on computational semigroup theory:

`compsemi.wordpress.com`

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Thank You!